

A double-contraction reductio at a degree-5 vertex via Kempe chains as Heawood face-chains

Status: exploratory strategy note. Records a proof skeleton and the single lemma it reduces to; nothing here is proved. Companion to `boundary-restriction-structure.tex` and `paper.tex`.

Setup: the move

Let G be a minimal counterexample to the Four Colour Theorem (a smallest triangulation with no proper 4-vertex-colouring). By Euler G has a vertex v of degree ≤ 5 ; degrees ≤ 4 are classically reducible, so take $\deg(v) = 5$. The link of v is a pentagon a, x, b, y, z (consecutive neighbours adjacent). For a non-adjacent (diagonal) pair, say a, b , the *double contraction* contracts both edges (a, v) and (b, v) , identifying $a = v = b$ into a single vertex. There are 5 diagonal pairs.

The double contraction is a proper minor with two fewer vertices, hence a smaller planar graph G' ; by minimality G' is 4-colourable. It stays a simple triangulation iff each contracted edge is non-separating (its endpoints have exactly two common neighbours).

Kempe chains as Heawood face-chains

Work in the cubic dual G' (vertices = faces of G , so the Heawood ± 1 labels sit on dual vertices). Tait-colour the edges with the three nonzero Klein-4 elements $\{a, b, c\}$ (edge colour = colour-difference across the primal edge). Then:

Observation (Kempe chain = Tait cycle = flip-set). *A primal Kempe chain (component of two vertex-colour classes of G) corresponds to a connected component of the two-edge-colour subgraph $\{a, b\}$ of the dual. That subgraph is 2-regular, so it is a disjoint union of cycles; each such cycle alternates a, b, a, b, \dots and is therefore **always even**. The cycle is literally a cyclic chain of faces of G . A Kempe swap reverses the rotational order (a, b, c) at exactly the vertices on the cycle, so the set of faces whose Heawood sign flips under the swap is precisely the Kempe chain.*

Observation (Sign alternation = same-side rule). *Even-ness does not force the Heawood signs to alternate along the cycle. At each cycle vertex the third (off-cycle, colour- c) edge points either inside or outside the region bounded by the cycle. For consecutive vertices A, B :*

same side (both in / both out) \Rightarrow signs flip; opposite side \Rightarrow signs repeat.

(Local chirality computation: the four in/out cases give $+, - / -, + / +, + / -, -$.) Hence the Heawood sign sequence around a Tait cycle is determined, up to one global sign, by the in/out pattern of the third edges; it is fully alternating iff there are zero side-switches, i.e. all third edges lie on one side.

The transport hypothesis (L1)

Conjecture (Chain transport, claimed for all triangulations). *The double contraction preserves the relevant (crossing) Kempe/Heawood chain structure as it propagates up the nested tire decomposition: the two chains anchored at the pentagon survive, consistently, through every tire interface above v .*

The reductio (two nested contradictions)

We are *not* claiming G' is uncolourable, and we are not claiming intertwined Kempe chains are impossible (they occur, e.g. in the Errera graph). The reducible object is *forced* intertwining.

Inner. Suppose intertwining is the *only* way to colour G' (every colouring is intertwined at v). By transport the two crossing chains pass through every interface above; if uncrossing is excluded at each interface, the tire restriction relations R_K collapse to their “crossed-only” sub-relations. Contract each interface **along the chain** to obtain a strictly smaller triangulation H'' . If the crossed-only sub-relations admit no compatible gluing on H'' , then H'' is uncolourable and smaller than G — contradicting minimality.

Outer. Therefore G' admits a non-intertwined colouring; it uncrosses at v , frees a colour for v , and lifts to a colouring of G — contradicting that G is a counterexample. \square (modulo the lemma below)

What it all reduces to

Claim (the only open content). *Under the crossed-only collapse, the surviving Heawood boundary sequences at a (forced-short) tire interface have empty pointwise-negation gluing on H'' — not merely small, genuinely empty.*

This is where the 2^{n-2} constraint floor and the mod-3 side-pattern must do real work: the floor guarantees that *smallness alone* never empties a relation, so the emptiness must come from the chains pinning the sub-relation down, not from short interface length.

Open points / to pin next

- **Define H'' precisely.** Current candidate: contract each tire interface along the transported chain (the global analogue of the local double contraction). Confirm this is a definite, strictly smaller triangulation.
- **Lift is not automatic.** “Non-intertwined” must be defined as “admits a swap collapsing the pentagon a, x, b, y, z to ≤ 3 colours,” or freeing v fails even after uncrossing (with $c(a) = c(b)$ the pentagon can still show 4 colours).
- **Errera oracle.** Run the whole construction against the Errera graph (and Fritsch / Kittell). These are colourable but Kempe-intertwined; the hypothesis “every colouring intertwined” must *fail* for them, so the construction must decline. *Why* it declines is exactly the mechanism the Claim must deny in the counterexample case.
- **Prove or break L1** as a statement about all triangulations.