

# Cut-and-depth-label: a procedure for labelling half-graphs of a 6-edge cut by “distance to the cut”

## Procedure

Given a maximal planar graph  $G$  and its dual  $G'$ :

1. Find a 6-edge cut  $C \subseteq E(G')$  partitioning  $V(G')$  into  $S$  and  $V \setminus S$  with both sides non-empty.
2. Remove the cut edges to obtain two graphs  $G'_0 := G'[S] \cup \emptyset_C$  and  $G'_1 := G'[V \setminus S] \cup \emptyset_C$  (each a cubic graph minus boundary edges).
3. In each  $G'_i$ :
  - (a) Let  $V_i$  be the set of vertices of degree 2 in  $G'_i$  (i.e. original cubic vertices incident to exactly one cut edge; the cut edge accounts for the missing edge, so the vertex sits at degree  $3 - 1 = 2$ ).
  - (b) For each  $v \in V_i$ , attach a new pendant edge to a fresh vertex, and **label these pendant edges with depth 0**.
  - (c) For  $d = 0, 1, 2, \dots$ : label every unlabelled edge that shares a vertex with a depth- $d$  edge with depth  $d + 1$ .
  - (d) Stop when every edge has a depth label.

**Interpretation.** The depth labels give a BFS distance in the line graph of  $G'_i$  starting from the pendants added at the cut. Equivalently, the depth of an edge  $e$  in  $G'_i$  is the minimum number of edges traversed (via shared-vertex adjacency) to reach a pendant.

**Caveat.** If the cut  $C$  is not a *matching* cut (i.e., the 6 cut edges share vertices), then some boundary vertices have degree  $< 2$  in  $G'_i$  and do not receive a pendant under the strict reading of step 3(a). When the cut is a matching, each of the 12 boundary vertices (6 per side) has degree exactly 2 and receives a pendant; the construction is symmetric.

## Example: Holton-McKay graph #0

We apply the procedure to the first of the six non-Hamiltonian 38-vertex cubic plane graphs found by Holton and McKay (loaded from `papers/even_level_graph_generators/experiments/nonham38m4.pc`). This  $G'$  is itself a cubic plane graph; its dual  $G$  is a 21-vertex triangulation.

**The chosen 6-edge cut.** Greedy search over 128 distinct 6-edge cuts in  $G'$ , preferring *matching cuts* (both sides have 6 distinct boundary vertices) and then balanced  $|S|$ , returns

$$|S| = 10, \quad C = \{(34, 29), (35, 30), (26, 22), (27, 23), (28, 24), (31, 25)\}.$$

This is a matching cut: the 6 edges have 12 distinct endpoints, 6 on each side.

## Resulting half-graphs.

	$G'_0$	$G'_1$
$ S $	10	28
Original vertices in $G'_i$	10	28
$ V_i $ (pendants added)	6	6
Total vertices in $G'_i$	16	34
Total edges	18	45
Max depth assigned	2	7

Multi-cut vertices (degree  $< 2$  in induced subgraph): *none* on either side, since the cut is a matching.

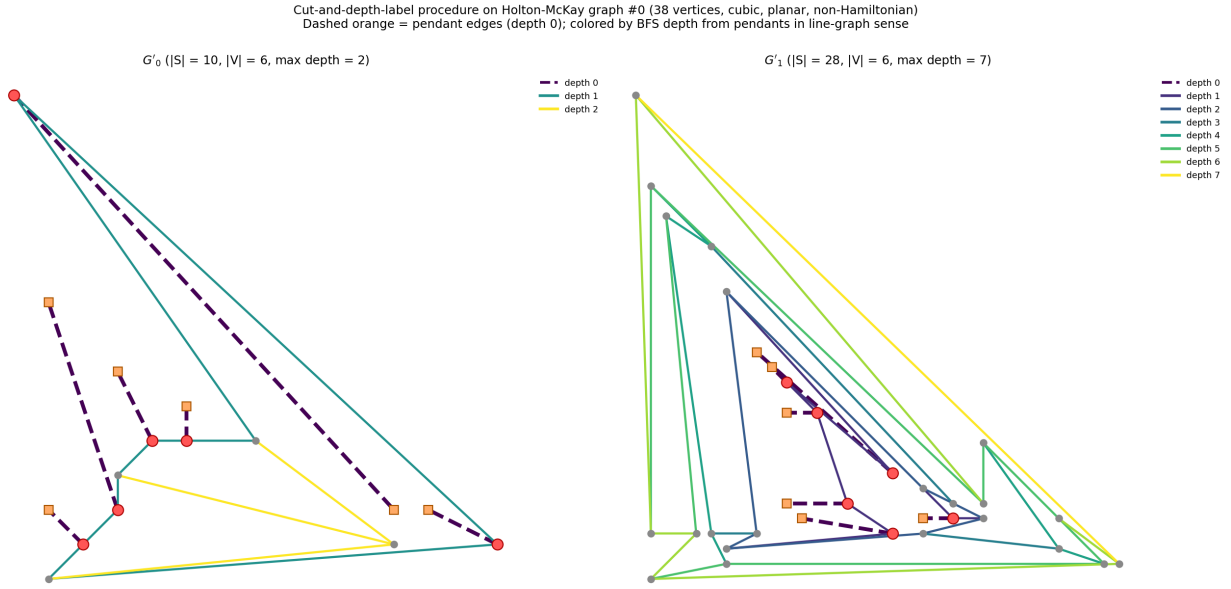


Figure 1: Cut-and-depth-label procedure on Holton-McKay graph #0 (38 vertices, 57 edges, cubic planar non-Hamiltonian). The 6-edge matching cut splits  $G'$  into a 10-vertex piece  $G'_0$  (left) and a 28-vertex piece  $G'_1$  (right). Pendant edges (dashed, dark purple) are at depth 0; the colour gradient encodes depth 0 through max depth via BFS in the line-graph sense. Orange squares are the new pendant vertices; red circles are the boundary vertices in  $V_i$ ; gray circles are interior vertices.

## Visualization.

## Connection to chain pigeonhole / 4CT reducibility

The procedure mirrors the 4CT cut-and-reglue scheme (`rainbow_proof.tex`, `worst_case_proof_sketch.tex`, `two_approaches_comparison.tex`) at the structural level. After cutting, each  $G'_i$  is a cubic-minus-boundary graph; the pendant additions formally restore cubicity at the 6 degree-2 boundary vertices. The depth label on each edge measures its “distance to the cut.”

For a minimum counterexample to the 4CT (i.e., a cubic plane  $G'$  with no proper 3-edge-colouring), the depth labels organise each  $G'_i$  into concentric layers indexed by distance to the cut. The 3-edge-colourings of  $G'_i$  must extend a colouring at the depth-0 pendants (= a ring colouring

at the cut); the BFS ordering by depth is the natural induction order for propagating the colouring inward.

In the tire framework, the cut cycle  $\gamma$  in the primal  $G$  (corresponding to the 6-edge cut in  $G'$  via planar duality) plays the role of the tire's inner boundary on one side and outer boundary on the other. The depth label on  $G'_i$ -edges is exactly the dual analogue of plane depth from  $\gamma$  (cf. the level-cycle generalization discussion in the recent conversation).

## Limitations of this example

- Holton-McKay graphs are cyclically 5-edge-connected (not 6-edge-connected), so 6-edge cuts are not the minimum cyclic cut. The smallest cyclic edge cuts in this graph are size 5.
- The matching 6-cut found is highly imbalanced ( $|S| = 10$  vs.  $|S^c| = 28$ ). Searching among the 128 distinct 6-edge cuts for a balanced matching cut may give better examples.
- Depth labels propagate via the line graph, not via vertex BFS. An alternative procedure would label *vertices* by BFS distance from the boundary; both yield similar layered structures but with slightly different counts.