

DUAL DECOMPOSITION OF MINIMAL COUNTEREXAMPLES

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ABSTRACT. We propose the *face-monochromatic-pair conjecture*, a structural property of proper 3-edge-colourings of cubic plane graphs that, if true, implies the Four Colour Theorem. Working in the planar dual G' of a hypothetical minimal counterexample G to 4CT, we delete a single pentagonal face of G' and rewire its five external vertices around a new apex vertex and a chord; the resulting *reduced dual* $\hat{G}'_{v,i}$ is a smaller cubic plane graph whose proper 3-edge-colourings, by the minimality of G , are constrained by a chord-apex condition and a pair of Kempe-cycle conditions. The face-monochromatic-pair conjecture, in its strengthened form, asserts the existence in every such colouring of a face F and two non-incident same-coloured edges $e_1, e_2 \in \partial F$ whose subdivision-and-bridging produces a 4-face f_n whose boundary colouring places it under the hypothesis of a 4-face edge-suppression theorem; we use this theorem to derive a proper 3-edge-colouring of G' , contradicting minimality. We verify the conjecture computationally on all chord-apex+Kempe colourings of reduced duals with $|V(G)| \leq 20$ (142,812 colourings, all pass); the weaker form is verified up to $|V(G)| \leq 21$ (535,182 colourings, all pass).

1. INTRODUCTION

The Four Colour Theorem (4CT) — that every loopless plane graph admits a proper 4-vertex-colouring — has, since the late 1970s, only been proved by computer-assisted case analysis on a discharging argument over a finite set of unavoidable reducible configurations. Appel and Haken’s original proof [1, 2], the Robertson–Sanders–Seymour–Thomas reworking [3], and Gonthier’s machine-checked version [4] all share that structure.

This paper takes a different approach: rather than discharge over configurations in the triangulation G , we work in its planar dual G' , a cubic plane graph whose proper 3-edge-colourings correspond by Tait’s theorem to proper 4-vertex-colourings of G . Assuming G is a minimal counterexample to 4CT, we delete a single pentagonal face of G' and rewire its five external vertices, obtaining a smaller cubic plane graph $\hat{G}'_{v,i}$ — the *reduced dual* — which by minimality is properly 3-edge-colourable. Two structural lemmas constrain every such colouring: a *chord-apex* condition (Lemma 3.6) forcing two named edges to share a colour, and a pair of *Kempe-cycle* conditions (Lemma 3.7) placing four of the rewired edges on common bichromatic Kempe cycles. These constraints are the starting point of the present development.

The main contribution of the paper is the *face-monochromatic-pair conjecture* (Conjecture 5.1) and its strengthening (Conjecture 5.3), which we show together imply the Four Colour Theorem. The supporting ingredients are the chord-apex

2010 *Mathematics Subject Classification*. Primary .

Key words and phrases. four colour theorem, plane triangulation, dual graph, cubic planar graph, edge connectivity, cyclic edge cut, Tait colouring, 3-edge-colouring.

and Kempe-cycle lemmas on reduced-dual colourings, the classical operation of *edge suppression* (delete the edge and smooth its two degree-2 endpoints; equivalently, simple-graph contraction in the dual triangulation; recalled in Section 4), and an observation that suppression preserves 3-edge-colourability when applied across a 4-face whose two opposite boundary edges carry different colours (Theorem 4.2).

The strategy is to construct, from a putative minimal counterexample’s reduced-dual colouring, a 4-face f_n in a slightly modified graph \widehat{G}'^+ to which the suppression theorem applies; the suppression then produces a properly 3-edge-coloured graph from which a 3-edge-colouring of G' can be recovered, contradicting the non-4-colourability of G . The face-monochromatic-pair conjecture asserts the existence of the structural data (F, e_1, e_2) needed to build f_n ; the strengthening guarantees that f_n ’s boundary colouring falls under the suppression theorem’s hypothesis. Both conjectures have been verified computationally on all chord-apex+Kempe colourings of reduced duals up to $|V(G)| \leq 20$, with the weaker form going up to $|V(G)| \leq 21$.

Organization. Section 2 fixes the minimal-counterexample framework: G is a triangulation, $\delta(G) \geq 5$, and every triangulation on fewer vertices is properly 4-colourable. Section 3 introduces the reduced dual $\widehat{G}'_{v,i}$ and proves the chord-apex and Kempe-cycle lemmas. Section 4 defines edge suppression and proves its 4-face 3-edge-colourability theorem. Section 5 states the two conjectures, reports the empirical verification, and gives the implication to 4CT.

Companion paper. An iterated version of the reduced-dual construction — producing a sequence H_1, H_2, \dots of progressively smaller cubic plane graphs and tracking an accumulating “protected” edge set — is the subject of a companion paper. The present paper does not use that iteration.

2. THE MINIMAL COUNTEREXAMPLE

Throughout, a *triangulation* is a simple plane graph, with a fixed embedding, in which every face — including the outer face — is bounded by a triangle. We first reduce to triangulations, then record the degree properties a smallest counterexample must have.

Lemma 2.1 (Reduction to triangulations). *If every triangulation is properly 4-vertex-colourable, then so is every plane graph.*

Proof. Let H be a plane graph. Add edges to H , maintaining planarity, until no further edge can be added; the result is a triangulation H^+ on the same vertex set with $E(H) \subseteq E(H^+)$. A proper 4-colouring of H^+ restricts to a proper 4-colouring of H , since every edge of H is an edge of H^+ . \square

By Lemma 2.1, if the Four Colour Theorem fails then it fails for some triangulation. We may therefore make the following assumption.

Definition 2.2 (Minimal counterexample). Let G be a triangulation on the fewest vertices that admits no proper 4-vertex-colouring. We call G a *minimal counterexample*. By minimality, every triangulation on fewer than $|V(G)|$ vertices is properly 4-colourable.

Remark 2.3. Since every triangulation on at most four vertices is properly 4-colourable (the largest being K_4), a minimal counterexample has $|V(G)| \geq 5$; the degree bound below sharpens this to $|V(G)| \geq 12$.

Lemma 2.4 (Minimum degree). *A minimal counterexample G has minimum degree $\delta(G) \geq 5$.*

Proof. Suppose some vertex v has $\deg(v) = d \leq 4$.

If $d \leq 3$, let $G' = G - v$. Then G' is a plane graph on fewer vertices, so by Definition 2.2 and Lemma 2.1 it has a proper 4-colouring. The at most three neighbours of v use at most three colours, so a fourth colour is free for v , extending the colouring to G — a contradiction.

If $d = 4$, again 4-colour $G - v$. If the four neighbours of v use at most three colours we extend as before, so assume they receive all four colours; let v_1, v_2, v_3, v_4 be the neighbours in cyclic order around v , coloured 1, 2, 3, 4. Consider the subgraph induced by the colour classes 1 and 3, and let K be its connected component containing v_1 . If $v_3 \notin K$, swap colours 1 and 3 on K ; now no neighbour of v is coloured 1, freeing it for v . If $v_3 \in K$, then a 1–3 Kempe chain joins v_1 to v_3 , and this chain together with v encloses exactly one of v_2, v_4 ; hence the 2–4 component containing v_2 cannot also reach v_4 , and swapping colours 2 and 4 on it frees colour 2 for v . Either way the colouring extends to G , a contradiction.

Hence $\delta(G) \geq 5$. □

3. THE REDUCED DUAL

Write G' for the dual of G : since G is a triangulation, G' is a cubic plane graph in which each vertex of G corresponds to a face of G' , each face of G to a vertex of G' , and each edge to a dual edge. A vertex of G of degree k corresponds to a k -gonal face of G' .

By Lemma 2.4, $\delta(G) \geq 5$, and Euler's formula gives $\sum_{u \in V(G)} (6 - \deg u) = 12$, so G has a vertex of degree exactly 5 (indeed at least twelve). Fix such a vertex v . Its dual face F_v is a pentagon, bounded by the five dual vertices corresponding to the five faces of G incident to v .

Definition 3.1 (Reduced dual). Let v be a degree-5 vertex of G with pentagonal dual face F_v , and fix an index $i \in \{0, 1, 2, 3, 4\}$. The *reduced dual* $\widehat{G}'_{v,i}$ is the plane graph obtained from G' as follows.

- (1) Delete the five dual vertices on the boundary of F_v , together with all edges incident to them. Each deleted vertex is cubic, with two edges on ∂F_v and one edge leaving F_v ; deleting the five boundary vertices therefore removes the five external edges as well, dropping their five outer endpoints from degree 3 to degree 2. These five degree-2 vertices lie on the boundary of a single face F of the resulting graph.
- (2) List the five degree-2 vertices in clockwise order around F as $A = (A_0, A_1, A_2, A_3, A_4)$.
- (3) Add a new vertex v_n and join it to A_i, A_{i+1} , and A_{i+2} (indices mod 5) by three new edges.
- (4) Add a new edge between A_{i+3} and A_{i+4} (indices mod 5).

Remark 3.2. Steps (3) and (4) restore cubicity: A_i, A_{i+1}, A_{i+2} each gain one edge to v_n and A_{i+3}, A_{i+4} each gain the new edge, so all five return to degree 3, and v_n has degree 3. Since A_i, \dots, A_{i+2} and A_{i+3}, A_{i+4} are each consecutive along ∂F , the new vertex and edge can be drawn inside F without crossings, so $\widehat{G}'_{v,i}$ is again a cubic plane graph. The construction depends on the choice of i up to the rotational symmetry of A .

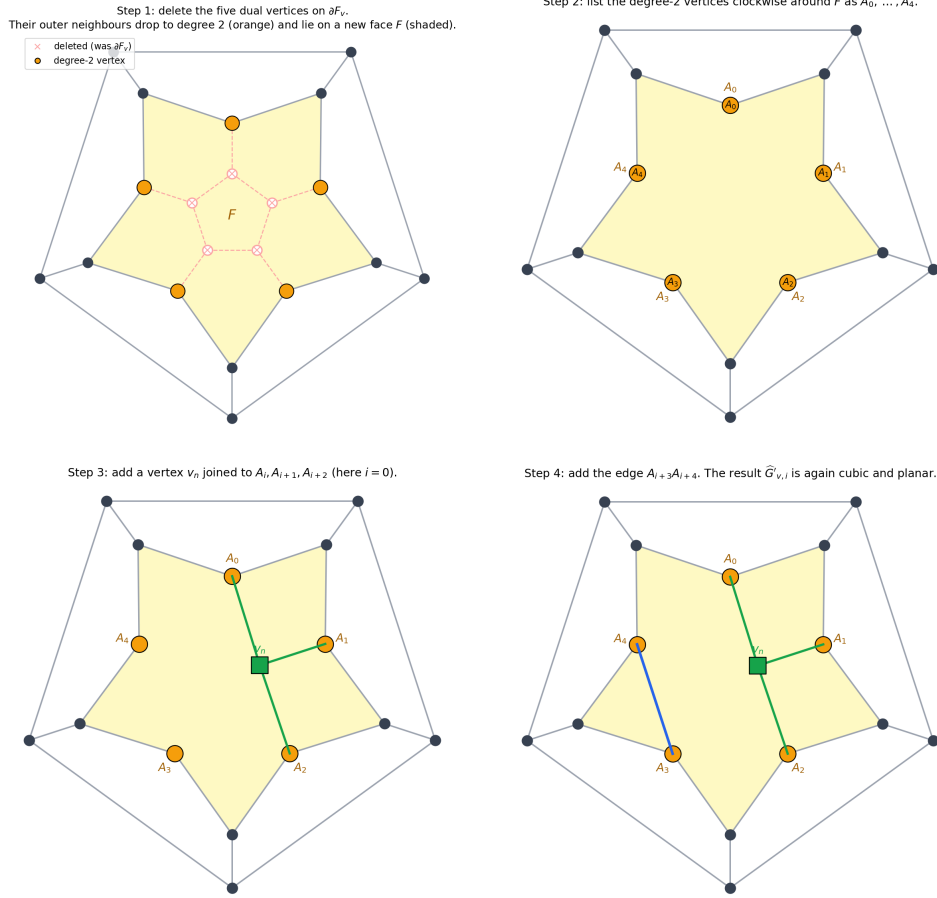


FIGURE 1. The four steps of Definition 3.1, illustrated on $G' =$ the dodecahedron (dual of the icosahedron) with F_v the inner pentagon and $i = 0$. Top left: delete the five boundary vertices of F_v , leaving five degree-2 vertices on a new face F . Top right: order them clockwise as A_0, \dots, A_4 . Bottom left: add v_n joined to A_0, A_1, A_2 . Bottom right: add the chord A_3A_4 , giving the cubic plane graph $\widehat{G}'_{v,0}$.

Definition 3.3 (Edges of the reduced dual). The four edges added in steps (3) and (4) of Definition 3.1 are named as follows. The chord $A_{i+3}A_{i+4}$ is the *merged edge*; the edge $A_{i+1}v_n$ is the *spike edge*; the edge A_iv_n is the *side-0 edge*; and the edge $A_{i+2}v_n$ is the *side-1 edge*. In the $i = 0$ case of Figure 1 these are $\{A_3, A_4\}$, $\{A_1, v_n\}$, $\{A_0, v_n\}$, and $\{A_2, v_n\}$ respectively.

We will use the following structural fact about proper 3-edge-colourings near a pentagonal face of a cubic plane graph; it is stated for a generic such graph H , not specifically for the reduced dual.

Lemma 3.4 (Pentagonal externals). *Let H be a cubic plane graph and F a pentagonal face of H , with ∂F traversed clockwise as u_0, u_1, u_2, u_3, u_4 . For each i let*

f_i be the unique edge of H incident to u_i that does not lie on ∂F . An assignment φ of colours from $\{1, 2, 3\}$ to the ten edges incident to $\{u_0, \dots, u_4\}$ is proper at every u_i if and only if there is some index j such that

$$\varphi(f_j) = \varphi(f_{j+1}) = \varphi(f_{j+2}) \quad \text{and} \quad \{\varphi(f_{j+3}), \varphi(f_{j+4})\} = \{1, 2, 3\} \setminus \{\varphi(f_j)\},$$

indices mod 5.

Proof. Write $e_i = u_i u_{i+1}$ for the boundary edges of ∂F (indices mod 5). A colouring φ is proper at every u_i if and only if at each u_i the three incident edges e_{i-1}, e_i, f_i receive three distinct colours; whenever this holds, $\varphi(f_i)$ is forced to be the unique colour in $\{1, 2, 3\} \setminus \{\varphi(e_{i-1}), \varphi(e_i)\}$, and φ restricts to a proper 3-edge-colouring of the cycle ∂F .

(\Rightarrow) The line graph of ∂F is C_5 , whose maximum independent set has size 2, so no colour appears more than twice on ∂F ; and since ∂F is an odd cycle, all three colours appear. The colour multiset on $(\varphi(e_0), \dots, \varphi(e_4))$ is therefore $(2, 2, 1)$, with the singleton at a unique position. Cyclically shifting indices we may place this position at 0; let c be the singleton colour. The remaining four edges form the path $e_1 e_2 e_3 e_4$, which by propriety alternates between the other two colours, so for some labelling $\{a, b, c\} = \{1, 2, 3\}$,

$$(\varphi(e_0), \varphi(e_1), \varphi(e_2), \varphi(e_3), \varphi(e_4)) = (c, a, b, a, b).$$

Reading off the forced values of $\varphi(f_i)$,

$$\varphi(f_0) = a, \quad \varphi(f_1) = b, \quad \varphi(f_2) = \varphi(f_3) = \varphi(f_4) = c,$$

which is the lemma's pattern at $j = 2$ (the cyclic shift maps this back to the corresponding j in the original indexing). This case is the unique proper 3-edge-colouring of ∂F up to cyclic shift and permutation of $\{1, 2, 3\}$ (since $5 \cdot 3! = 30 = P(C_5, 3)$, the chromatic polynomial of C_5 at 3), so it exhausts every proper φ .

(\Leftarrow) The lemma's hypothesis is invariant under cyclic shifts of indices and under permutations of $\{1, 2, 3\}$, so we may assume $j = 2$, $\varphi(f_2) = \varphi(f_3) = \varphi(f_4) = c$, $\varphi(f_0) = a$, and $\varphi(f_1) = b$, with $\{a, b, c\} = \{1, 2, 3\}$. Propriety at u_i and u_{i+1} requires $\varphi(e_i) \notin \{\varphi(f_i), \varphi(f_{i+1})\}$, which gives

$$\varphi(e_0) = c, \quad \varphi(e_1) = a, \quad \varphi(e_2) \in \{a, b\}, \quad \varphi(e_3) \in \{a, b\}, \quad \varphi(e_4) = b.$$

The remaining propriety condition $\varphi(e_{i-1}) \neq \varphi(e_i)$ holds automatically at u_0, u_1, u_4 , forces $\varphi(e_2) = b$ at u_2 , and then forces $\varphi(e_3) = a$ at u_3 . The resulting triples $(\varphi(e_{i-1}), \varphi(e_i), \varphi(f_i))$ at u_0, u_1, u_2, u_3, u_4 are

$$(b, c, a), \quad (c, a, b), \quad (a, b, c), \quad (b, a, c), \quad (a, b, c),$$

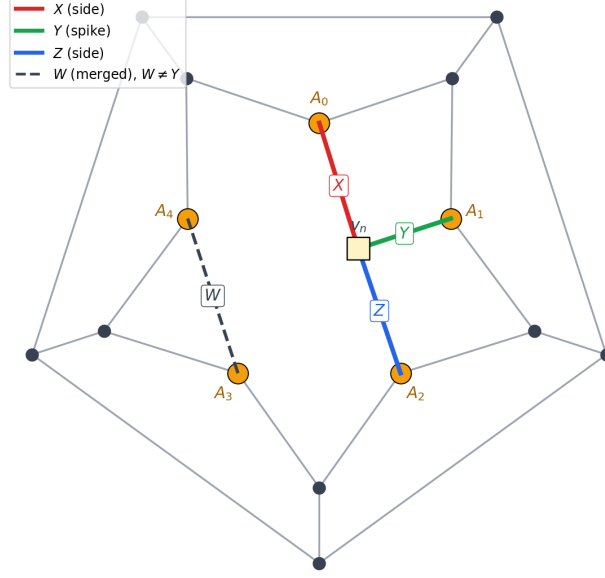
each a permutation of $\{1, 2, 3\}$, so φ is proper at every u_i . \square

Remark 3.5. The two-element condition $\{\varphi(f_{j+3}), \varphi(f_{j+4})\} = \{1, 2, 3\} \setminus \{\varphi(f_j)\}$ cannot be dropped: a 3-colouring satisfying $\varphi(f_j) = \varphi(f_{j+1}) = \varphi(f_{j+2})$ alone need not extend, e.g. $(1, 1, 1, 1, 2)$.

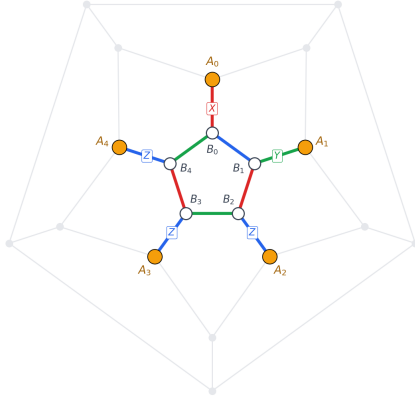
Since $\widehat{G}'_{v,i}$ is the dual of a triangulation on fewer vertices than G , it is 3-edge-colourable by the minimality of G . The following lemma constrains every such colouring.

Lemma 3.6. *Let G be a minimal counterexample, and let $\widehat{G}'_{v,i}$ be a reduced dual of its dual G' . Then in every proper 3-edge-colouring of $\widehat{G}'_{v,i}$, the merged edge and the spike edge receive the same colour.*

Step 1: φ on $\widehat{G}_{v,0}$ assigns distinct colours X, Y, Z to the v_n -edges (propriety at v_n);
by hypothesis $W \neq Y$, forcing $W \in \{X, Z\}$.



Step 2: lift to G' when $W = Z$. The externals inherit $\psi(f) = (X, Y, Z, Z, X)$;
Lemma 2.4 colours the five edges of ∂F_v .



Step 3: lift to G' when $W = X$. The externals inherit $\psi(f) = (X, Y, Z, X, X)$;
Lemma 2.4 colours the five edges of ∂F_v .

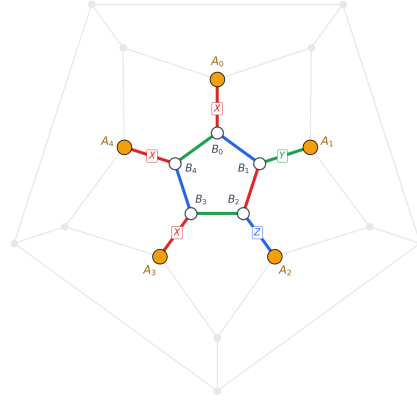


FIGURE 2. The proof of Lemma 3.6, illustrated for $i = 0$ on $G' =$ the dodecahedron. Top: under the assumption $W \neq Y$, propriety at v_n forces $W \in \{X, Z\}$. Bottom: in either case the lift to G' has externals satisfying the hypothesis of Lemma 3.4, which colours ∂F_v to extend ψ to a proper 3-edge-colouring of G' .

Proof. After cyclically relabelling, assume $i = 0$. Suppose for contradiction that φ is a proper 3-edge-colouring of $\widehat{G}_{v,0}'$ in which the merged edge $\{A_3, A_4\}$ and the spike edge $\{A_1, v_n\}$ receive different colours (Figure 2, top), and write

$$X = \varphi(\{A_0, v_n\}), \quad Y = \varphi(\{A_1, v_n\}), \quad Z = \varphi(\{A_2, v_n\}), \quad W = \varphi(\{A_3, A_4\}).$$

Propriety of φ at v_n forces $\{X, Y, Z\} = \{1, 2, 3\}$, and the assumption $W \neq Y$ leaves $W \in \{X, Z\}$.

We lift φ to a colouring ψ of $E(G')$ as follows. Let B_0, \dots, B_4 be the boundary vertices of ∂F_v in G' , indexed so that $f_k = B_k A_k$. On every edge of G' that survived the reduction, set $\psi = \varphi$. At each A_k the two surviving edges retain their φ -colours, so the remaining edge at A_k — in G' this is the external f_k ; in $\widehat{G}'_{v,0}$ this is a v_n -edge ($k \in \{0, 1, 2\}$) or the chord ($k \in \{3, 4\}$) — is forced to take the third colour at A_k . Since the two-surviving-edge colours at A_k agree in G' and $\widehat{G}'_{v,0}$, the third colour does too, giving

$$\psi(f_0) = X, \quad \psi(f_1) = Y, \quad \psi(f_2) = Z, \quad \psi(f_3) = \psi(f_4) = W$$

(the last two equalities holding because the chord is a single edge contributing its colour at each of A_3 and A_4).

It remains to assign colours to the five boundary edges $B_k B_{k+1}$ of ∂F_v . Apply Lemma 3.4 to G' at the face F_v with the B_k 's as its boundary vertices and the same indexing. The external vector $(\psi(f_0), \dots, \psi(f_4)) = (X, Y, Z, W, W)$ falls into one of two cases (Figure 2, bottom):

- if $W = Z$, it is (X, Y, Z, Z, Z) : three consecutive Z 's at positions 2, 3, 4, with $\{X, Y\} = \{1, 2, 3\} \setminus \{Z\}$;
- if $W = X$, it is (X, Y, Z, X, X) : three consecutive X 's at positions 3, 4, 0, with $\{Y, Z\} = \{1, 2, 3\} \setminus \{X\}$.

Each case satisfies the hypothesis of Lemma 3.4; its (\Leftarrow) direction therefore assigns colours to the boundary edges $B_k B_{k+1}$ that make ψ proper at every B_k .

The resulting ψ is a proper 3-edge-colouring of G' : proper at every B_k by the lemma, at every A_k by the construction, and at every other vertex because such a vertex has the same neighbourhood in G' as in $\widehat{G}'_{v,0}$ with the same incident-edge colours. By Tait's theorem, G' is 3-edge-colourable iff G is 4-vertex-colourable, contradicting that G is a counterexample. The assumption $W \neq Y$ is therefore false. \square

For a pair of colours $\{a, b\} \subseteq \{1, 2, 3\}$, the subgraph of $\widehat{G}'_{v,i}$ on the edges coloured a or b is 2-regular (since at each vertex exactly one of the three incident edges is excluded), and hence a disjoint union of cycles. We call each such cycle a $\{a, b\}$ -Kempe cycle, and reserve the notation for the specific cycle containing a given edge when the context makes it clear. Swapping the two colours on a single Kempe cycle yields another proper 3-edge-colouring of the same graph.

Lemma 3.7 (Kempe cycles through the spike). *Let G be a minimal counterexample, fix a reduced dual $\widehat{G}'_{v,i}$ of G' , and let φ be a proper 3-edge-colouring of $\widehat{G}'_{v,i}$. Write c for the common colour assigned by φ to the spike and the merged edge (Lemma 3.6), and c_0, c_1 for the colours of the side-0 and side-1 edges respectively, so $\{c, c_0, c_1\} = \{1, 2, 3\}$. Then*

- (1) *the $\{c, c_0\}$ -Kempe cycle through the spike edge contains both the side-0 edge and the merged edge;*
- (2) *the $\{c, c_1\}$ -Kempe cycle through the spike edge contains both the side-1 edge and the merged edge.*

Proof. We prove (1); (2) is the same argument with c_1 and the side-1 edge in place of c_0 and the side-0 edge.

The spike edge $\{A_{i+1}, v_n\}$ and the side-0 edge $\{A_i, v_n\}$ share the vertex v_n and receive the two colours c, c_0 , so they both lie on the $\{c, c_0\}$ -Kempe cycle through v_n . Suppose for contradiction that the merged edge lies on a different $\{c, c_0\}$ -Kempe cycle K (it lies on *some* such cycle, since it has colour c). Let φ' be obtained from φ by swapping the colours c and c_0 along K alone: this is a Kempe swap, so φ' is again a proper 3-edge-colouring of $\widehat{G}'_{v,i}$. Under φ' the spike edge — which is not on K — still has colour c , but the merged edge — which is on K — now has colour c_0 . Hence in φ' the spike and the merged edge receive distinct colours, contradicting Lemma 3.6 applied to φ' . \square

4. EDGE SUPPRESSION

We recall the classical operation of *edge suppression* on cubic plane graphs: delete the edge and smooth the two resulting degree-2 endpoints. Under planar duality this coincides with simple-graph contraction on the dual side. It will be the central tool in Section 5 below, where we formulate a sufficient condition for the Four Colour Theorem.

Definition 4.1 (Edge suppression). Let H be a cubic plane graph and $e = uv$ an edge of H with $u \neq v$ and no edge of H parallel to e . The *edge suppression* of H at e is the graph H' obtained in two steps:

- (1) *Delete* the edge e ; the endpoints u and v each drop to degree 2.
- (2) *Smooth* each of u and v : at u , replace u and its two remaining incident edges ua, ub by a single new edge ab ; do the same at v . Both vertices u and v are removed, and two new edges are added in their place.

Provided the smoothings do not introduce a loop or parallel edge, H' is again a cubic plane graph, with $|V(H')| = |V(H)| - 2$ and $|E(H')| = |E(H)| - 3$.

Equivalently, H' is the planar dual of $\text{dual}(H)/e^*$, where e^* is the edge of $\text{dual}(H)$ crossing e and the contraction on the right-hand side is simple-graph contraction (loops removed, parallel edges absorbed). Under planar duality, contracting e^* in $\text{dual}(H)$ merges the two triangular faces of $\text{dual}(H)$ incident to e^* , and the parallel-edge cleanup corresponds exactly to the smoothing step on the primal side.

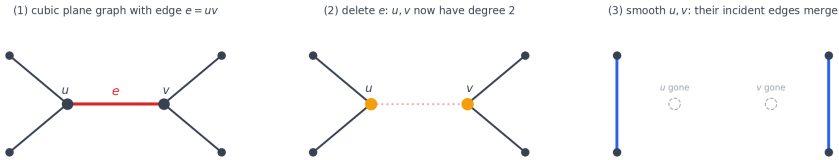


FIGURE 3. Edge suppression (Definition 4.1). Left: a fragment of a cubic plane graph with the suppressed edge $e = uv$ highlighted in red. Middle: deleting e leaves u and v of degree 2. Right: smoothing u and v replaces each pair of incident edges by a single new edge, removing u, v and giving a cubic plane graph again.

Theorem 4.2 (Edge suppression across a 4-face preserves 3-edge-colourability). *Let H be a cubic plane graph with a proper 3-edge-colouring φ , let f be a face of H with $|\partial f| = 4$, and let e_0, e_1 be the two edges of ∂f sharing no endpoint (the*

opposite pair on the 4-cycle ∂f). If $\varphi(e_0) \neq \varphi(e_1)$ and the edge suppression of H at e_0 (Definition 4.1) is well-defined (no loops or parallel edges are created), then the suppressed graph admits a proper 3-edge-colouring.

Proof. Write ∂f as the 4-cycle $v_0v_1v_2v_3$ with $e_0 = v_0v_1$ and $e_1 = v_2v_3$ (so e_0, e_1 are opposite); the remaining two boundary edges of f are $e_2 := v_1v_2$ and $e_3 := v_3v_0$. Since H is cubic, each v_i has exactly one edge not on ∂f : write w_i for that edge and u_i for its other endpoint, so $w_i = v_iu_i$ with $u_i \notin \{v_0, v_1, v_2, v_3\}$, for each $i \in \{0, 1, 2, 3\}$. Put $a := \varphi(e_0)$, $b := \varphi(e_1)$, and let c be the third colour.

Forced colours on the face. Propriety at v_1 and v_2 forces $\varphi(e_2) \notin \{a, b\}$, so $\varphi(e_2) = c$; then $\varphi(w_1) = b$ and $\varphi(w_2) = a$. Symmetrically $\varphi(e_3) = c$, $\varphi(w_0) = b$, and $\varphi(w_3) = a$. In particular $\varphi(w_0) = \varphi(w_1) = b$.

Construction of φ' . Let H' denote the edge suppression of H at e_0 ; its new edges are $e'_3 := v_3u_0$ (replacing e_3 and w_0 via the smoothing at v_0) and $e'_2 := v_2u_1$ (replacing e_2 and w_1 via the smoothing at v_1). Define $\varphi': E(H') \rightarrow \{1, 2, 3\}$ by

$$\varphi'(e) := \begin{cases} c & \text{if } e = e_1, \\ b & \text{if } e \in \{e'_2, e'_3\}, \\ \varphi(e) & \text{otherwise.} \end{cases}$$

That is: give each smoothed-in edge the colour b (the colour of the two w_i it absorbs), recolour e_1 to c , and leave every other edge of H' with its φ -colour.

Propriety. Every vertex of H' other than v_2, v_3, u_0, u_1 has the same incident edges and the same φ' -colours as it did under φ , so propriety is inherited there. At the four affected vertices,

vertex	edges in H'	colours under φ'
v_2	e_1, w_2, e'_2	c, a, b
v_3	e_1, w_3, e'_3	c, a, b
u_0	e'_3, α_0, β_0	b, a, c
u_1	e'_2, α_1, β_1	b, a, c

where α_i, β_i are the two edges of H at u_i other than w_i , whose φ -colours are forced to $\{a, c\}$ by propriety at u_i (since $\varphi(w_i) = b$). Each row lists three distinct colours, so φ' is proper. \square

5. THE FACE-MONOCROMATIC-PAIR CONJECTURE AND THE FOUR COLOUR THEOREM

The following conjecture identifies a structural property of every proper 3-edge-colouring of a reduced dual of a minimal counterexample. If true, it implies the Four Colour Theorem via Theorem 4.2.

Conjecture 5.1 (Face-monochromatic-pair conjecture). *Let G be a minimal counterexample to the Four Colour Theorem, and let $\hat{G}'_{v,i}$ be a reduced dual of $G' = \text{dual}(G)$. Then for every proper 3-edge-colouring φ of $\hat{G}'_{v,i}$ there exist a face F of $\hat{G}'_{v,i}$ and two distinct edges $e_1, e_2 \in \partial F$, with neither e_1 nor e_2 equal to the merged edge, such that:*

- (1) $\varphi(e_1) = \varphi(e_2)$. Write $a := \varphi(e_1) = \varphi(e_2)$.
- (2) e_1, e_2 , and the merged edge all lie on a common $\{a, b\}$ -Kempe cycle of φ , for some colour $b \neq a$.

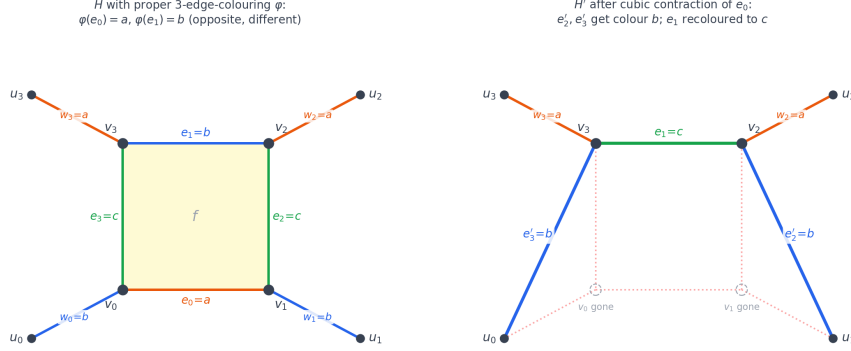


FIGURE 4. The recolouring used in the proof of Theorem 4.2. Left: the 4-face f of H under φ , with the forced colours $\varphi(e_0) = a$, $\varphi(e_1) = b$, $\varphi(e_2) = \varphi(e_3) = c$, $\varphi(w_0) = \varphi(w_1) = b$, and $\varphi(w_2) = \varphi(w_3) = a$. Right: the suppressed graph H' under φ' . The smoothed-in edges e'_2, e'_3 inherit the colour b from w_0, w_1 , and e_1 is recoloured from b to c ; every edge outside the face neighbourhood keeps its φ -colour (dotted in red: the five edges of H removed by the suppression).

- (3) *Exactly one edge of ∂F lies between e_1 and e_2 along one of the two arcs of ∂F . Equivalently, subdividing e_1, e_2 by new vertices X_1, X_2 and joining them by a new edge X_1X_2 inside F creates a new face f_n bounded by exactly 4 edges (the new edge X_1X_2 , the two subdivision halves adjacent to it, and the single ∂F -edge between e_1 and e_2).*

Remark 5.2. The conjecture cannot be tested on actual minimal counterexamples (none exist by the Four Colour Theorem), but its conclusion is checkable on the structural surrogates: proper 3-edge-colourings of reduced duals that satisfy both the chord-apex condition (Lemma 3.6) and the Kempe-cycle conditions (Lemma 3.7), since a minimal counterexample's reduced dual is forced to admit such colourings. For every min-degree-5 triangulation G with $|V(G)| \leq 21$, every pentagonal face F of G' , and every reduction index $i \in \{0, \dots, 4\}$, we enumerated all such colourings and verified the three clauses of Conjecture 5.1 (see `experiments/check_conj_face_kempe_scaled.py`); $n = 22$ ran past a 1800s budget after 641,700 colourings (all pass) without finishing the full set of 651 triangulations.

n	#tri	#col. tested	#sat.	status
12	1	0	—	vacuous (icosahedron)
13	0	—	—	no min-deg-5 tri
14	1	216	216	all pass
15	1	0	—	vacuous
16	3	864	864	all pass
17	4	4,650	4,650	all pass
18	12	8,070	8,070	all pass
19	23	21,138	21,138	all pass
20	73	107,874	107,874	all pass
21	192	392,370	392,370	all pass
22 (part.)	651	641,700	641,700	timeout
total ($n \leq 21$)	311	535,182	535,182	

The vacuous rows ($n = 12, 15$) are those where the relevant reduced duals admit no chord-apex+Kempe colourings, so the conjecture has no content there.

The following strengthening adds a fourth clause that arranges the new 4-face f_n to satisfy the hypotheses of Theorem 4.2.

Conjecture 5.3 (Strengthening of Conjecture 5.1). *Let G , $\widehat{G}'_{v,i}$, φ be as in Conjecture 5.1. Then there exist F , e_1 , e_2 satisfying clauses (1)–(3) of that conjecture, together with the following additional clause.*

- (4) *Let \widehat{G}'^{+} be the modified graph obtained from $\widehat{G}'_{v,i}$ by the modification of clause (3), and let φ' be the proper 3-edge-colouring of \widehat{G}'^{+} obtained from φ by swapping the two colours along the (subdivided) $\{a, b\}$ -Kempe cycle of clause (2) and assigning the new edge X_1X_2 the remaining (third) colour c . (Equivalently: φ' agrees with φ on every edge of $\widehat{G}'_{v,i}$ outside that Kempe cycle, and at X_1, X_2 the two subdivision halves take the colours $\{a, b\}$ in the order forced by propriety.) Then either*
- (i) *∂f_n uses all three colours under φ' , or*
 - (ii) *the $\{b, c\}$ -Kempe cycle of φ' through X_1X_2 is incident to exactly one edge of ∂f_n (namely X_1X_2 itself).*

Remark 5.4. The strengthened conjecture was tested on the same chord-apex+Kempe colourings as Remark 5.2; for each colouring we sought any Conjecture-5.1-witness (F, e_1, e_2) whose accompanying f_n satisfies clause (4) (see `experiments/check_conj_3.8_scaled.py`):

n	#tri	#col. tested	#sat.	status
14	1	216	216	all pass
16	3	864	864	all pass
17	4	4,650	4,650	all pass
18	12	8,070	8,070	all pass
19	23	21,138	21,138	all pass
20	73	107,874	107,874	all pass
total ($n \leq 20$)	116	142,812	142,812	

A subtlety: only about half of the Conjecture-5.1-witnesses individually satisfy clause (4) on each colouring, but in every case *some* witness does. Clause (4)

is therefore an existential statement at the witness level, not a property of every witness.

Remark 5.5 (The implication to the Four Colour Theorem). Clause (4)(i) of Conjecture 5.3 says that ∂f_n uses all three colours under φ' . Because ∂f_n is a 4-cycle and adjacent edges of \widehat{G}'^+ carry distinct φ' -colours, the cyclic colour pattern on ∂f_n must be (c, a, c, b) up to rotation and relabelling, with the two c -edges opposite and the two remaining opposite edges carrying the distinct colours a and b . Those two opposite edges therefore satisfy the hypothesis of Theorem 4.2: they lie on the 4-face f_n , share no endpoint, and have different φ' -colours. Theorem 4.2 then produces a proper 3-edge-colouring of the edge suppression of \widehat{G}'^+ at the a -coloured one.

Case (ii) of clause (4) is conjecturally reducible to case (i) by a single Kempe swap on the $\{b, c\}$ -cycle through X_1X_2 : by hypothesis that cycle is incident to ∂f_n only at X_1X_2 , so the swap flips $\varphi'(X_1X_2)$ from c to b while leaving the other three edges of ∂f_n unchanged — placing ∂f_n into the three-colour pattern of case (i).

Consequence. Theorem 4.2 now produces a proper 3-edge-colouring of the edge suppression of \widehat{G}'^+ at the chosen edge of f_n . Combined with the chord-apex and Kempe-cycle structure of $\widehat{G}'_{v,i}$ (Lemmas 3.6 and 3.7), this yields a proper 3-edge-colouring of G' , and by Tait's correspondence a proper 4-vertex-colouring of G — contradicting the assumption that G is a minimal counterexample. Hence Conjecture 5.3 implies the Four Colour Theorem.

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